## Lecture 2

Triangle Mesh and its Data Structure

## Overview

- The data structure to represent surfaces $\rightarrow$ efficiency and memory consumption of the geometric modeling algorithms
- We will
- Give a brief overview of the various data structures for mesh representations in the literature
- Elaborate the half-edge data structure, which is commonly used in modeling/processing 3D data


## Surfaces defined by triangle meshes

- A surface (2-manifold, two-dimensional manifold): a continuous topological space with infinitely many points, where each point has a local neighborhood homeomorphic to a 2-dimensional Euclidean space $\mathrm{E}^{2}$
- A triangle mesh is its approximation:
- We use a finite number of vertices and triangles
- Simply a collection of these triangles without any mathematical structure
- But it defines a piecewise linear representation with quadratic approximation


## Piecewise Linear Representation by Barycentric Parameterization

- Every triangle $T$ is determined by its three vertices $\mathbf{v}_{1}, \mathbf{v}_{2}, \mathbf{v}_{3}$, we denote it as $\mathrm{T}=\left[\mathrm{v}_{1}, \mathbf{v}_{2}, \mathbf{v}_{3}\right]$
- Any point $p$ in the interior of $T$ can be represented uniquely using a barycentric combination of the corner points:

$$
\mathbf{p}=a \mathbf{v}_{1}+b \mathbf{v}_{2}+c \mathbf{v}_{3}
$$

with $a+b+c=1, a, b, c>=0$.
Therefore, based on this per-triangle mapping, a 2D parameterization can be defined $\mathrm{f}: \boldsymbol{R}^{2} \rightarrow \boldsymbol{R}^{3}$, to represent the entire continuous surface approximated by this triangle mesh (details and algorithms will be discussed later, on how to actually construct this 2D layout $\rightarrow$ parameterization)

## Smooth Surfaces can be Well Approximated

- A sufficiently smooth surface is approximated by a triangle mesh (piecewise linear function)
- approximation error $\mathrm{O}\left(\mathrm{h}^{2}\right)$, with $\mathrm{h}=$ maximum edge length
- Error Reduced by a factor of $\frac{1}{4}$ : if we evenly split all edges
- A simple subdivision scheme for this: cutting each triangle into 4
- Face number: $\mathrm{F} \rightarrow 4 \mathrm{~F}$
- Approximation error : inversely proportional to F
- The actual approximation error depends on the $2^{\text {nd }}$ order Taylor expansion, i.e., on the curvature of the underlying smooth surface
- But roughly: a sufficient approximation can be obtained using moderate mesh complexity (you may want to adaptively adjust vertex density according to surface curvature, will be discussed later in meshing sessions)



## Geometric and Topological Components of a Triangle Mesh

- A triangle mesh has two components:
- Geometric components: vertex table $\rightarrow$ positions of points
- Topological components: face table $\rightarrow$ the graph encoding the connectivity
- What if you fix the topology (face table), and change the positions of vertices?
$\rightarrow$ a continuously deforming surface
- What if you only have the positions of sampled points, but don't know the connectivity?
- $\rightarrow$ could be complicated... different connectivity indicates different shapes


## Face-based Data Structure

- A simplest way to represent a surface mesh
- Storing a set of faces represented by their vertex positions
- also called "triangle soup"
- used in the stereolithography (STL) format
- if using $x$ (e.g. 32) bits to represent a vertex coordinate
- Each triangle needs $3 * 3 * x / 8=36$ bytes
- No connectivity info stored
- Inefficient for many geometric computing: e.g. traversing local adjacency information
- Vertex positions replicated as many times as the degree of the vertices

| Triangles |  |  |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $\mathbf{x}_{11}$ | Y11 | $\mathbf{Z}_{11}$ | $\mathbf{x}_{12}$ | $\mathrm{Y}_{12}$ | $\mathrm{Z}_{12}$ | $\mathrm{X}_{13}$ | $\mathrm{Y}_{13}$ | $\mathbf{Z}_{13}$ |
| $\mathrm{X}_{21}$ | $Y_{21}$ | $\mathbf{Z}_{21}$ | $\mathrm{X}_{22}$ | $Y_{22}$ | $\mathbf{Z}_{22}$ | $\mathrm{X}_{23}$ | $Y_{23}$ | $\mathrm{Z}_{23}$ |
|  | -•• |  | -•• |  |  | -•• |  |  |
| -•• |  |  | -•• |  |  | -•• |  |  |
| -•• |  |  | -•• |  |  | -•• |  |  |
| $\mathrm{X}_{\mathrm{F} 1}$ | YF1 | $\mathrm{Z}_{\mathrm{F} 1}$ | $\mathrm{X}_{\mathrm{F} 2}$ | $Y_{\text {F2 }}$ | $\mathrm{Z}_{\mathrm{F} 2}$ | $\mathrm{X}_{\mathrm{F} 3}$ | YF3 | $\mathrm{Z}_{\mathrm{F} 3}$ |

## Face-based Data Structure (2)

- An improved face-based data structure:
- To prevent the redundancy by indexed face set
- Stores an array of vertices
- Stores faces as sets of indices into this array
- Simple and efficient in storage
- Widely used in many formats such as OFF, OBJ, VRML, as well as our .M files
- if using x (e.g. 32) bits to represent a vertex coordinate and face indices
- Each vertex requires $3^{*} x / 8=12$ bytes
- Each triangle needs $3 * x / 8=12$ bytes
- (Roughly F=2V, by Euler formula)
- So on average: 18 bytes / triangle
- Only a half storage space
- No connectivity info stored
- Inefficient for many geometric computing: e.g. traversing local adjacency information

| Vertices | Triangles |
| :---: | :---: |
| $\mathrm{x}_{1} \quad \mathrm{y}_{1} \quad \mathrm{z}_{1}$ | $\mathrm{i}_{11} \quad i_{12} \quad i_{13}$ |
| $\ldots$ | $\ldots$ |
| $\mathrm{x}_{\mathrm{V}} \quad \mathrm{y}_{\mathrm{V}} \quad \mathrm{z}_{\mathrm{V}}$ | $\ldots$ |
|  | $\ldots$ |
|  | $\ldots$ |
| $\mathbf{i}_{\mathrm{F} 1}$ | $i_{\mathrm{F} 2}$ |

## What Does Geometric Computing Need?

- Access to individual elements (vertices edges, and faces): enumeration of all elements
- Local traversal, e.g.:
- What are the edges in a given face;
- What are the vertices in a given face or edge;
- What are incident faces of a given edge;
- What are incident faces or edges of a given vertex; ...
- Can you develop efficient algorithms to do these using the previous face-based data structure?
> An improved face-based data structure for efficient local traversal:
- For each face: store references to its 3 vertices + references to its neighboring triangles
- For each vertex: store a reference to its neighboring triangle +3 coordinates
- Used in CGAL for representing 2D Triangulation, 32 bytes / triangle (google CGAL)
- Enumerating the one-ring of a vertex is not easy
- Not easily extendable to general/mixed polygonal meshes



## Edge-based Data Structure

- A more generally used data structure, since the connectivity is a graph, directly relates to the mesh edges
- Many well known methods: winged-edge [Baumgart 72], quadedge [Guibas and Stolfi 85], and variants [O'Rourke 94]
- An example: Winged-edge structure
- Each edge stores references to its endpoint vertices + two incident faces + next and previous edge within the left and right faces
- Each vertex stores a reference to one of its incident edges
- Each face stores a reference to one of its incident edges
- 60 bytes / triangle
- Still not easy to traversing the onering (e.g. to traverse the one-ring of a vertex $v$, how do you know if it is the first or second vertex of an edge?)


| Edge |  |
| :--- | :--- |
| VertexRef | vertex[2] |
| FaceRef | face[2] |
| EdgeRef | next[2] |
| EdgeRef | prev[2] |

## Half-Edge Data Structure

- (What?) A common way to represent triangular mesh for geometric processing
- We first focus on triangle-mesh, (it works for general polygonal mesh).
- 3D analogy: half-face data structure for tetrahedral mesh
- (Why?) Effective for maintaining incidence information of vertices
- Efficient local traversal
- Relatively low spatial cost
- Supporting dynamic local updates/manipulations (edge collapse, vertex split, etc.)
- (Resources?) Codes are provided on the course website. After the class, please go through them carefully, we will work on it during the whole semester.


## Half-Edge Data Structure (cont.)

02 vertices share an edge, 2 faces share an edge
-Each face has 3 vertices,
$\square \rightarrow$ To store all adjacency information on half-edges
$\square$ Each edge has 2 half-edges (the boundary edge has 1)


## Half-Edge Data Structure (cont.)

DHalfedges are oriented consistently in counterclockwise order around each face
-Each halfedge designates a unique corner on each face (can be used to store texture coordinates, later in texture mapping)

- For each halfedge, we store:
- the vertex it points to (its target);
- its adjacent face (the face this halfedge locates);
- the next halfedge of the face;
- the previous halfedge in the face;
- its twin halfedge;
- For each vertex: store one of its incident incoming
 halfedges
- For each face: store one of its halfedges
- For each edge: store its two halfedges

D\# of halfedges H is about 6 times of V :
$\square \rightarrow 72$ bytes/triangle

## Half-Edge Data Structure (example)

1). Containers store primitives:

| The Vertex Container ${ }^{*}$ | $\mathrm{v} 1 \ldots \mathrm{v} 6$ |
| :--- | :--- |
| The Half-Edge Container | $[\mathrm{v} 1, \mathrm{v} 2],[\mathrm{v} 2, \mathrm{v} 3],[\mathrm{v} 3, \mathrm{v} 1],[\mathrm{v} 1, \mathrm{v} 3],[\mathrm{v} 3, \mathrm{v} 4], \ldots$ |
| The Edge Container | $[\mathrm{v} 1, \mathrm{v} 3],[\mathrm{v} 1, \mathrm{v} 2],[\mathrm{v} 2, \mathrm{v} 3],[\mathrm{v} 1, \mathrm{v} 4],[\mathrm{v} 3, \mathrm{v} 4], \ldots$ |
| The Face Container | $\mathrm{f} 1[\mathrm{v} 1, \mathrm{v} 2, \mathrm{v} 3] \ldots \mathrm{f}[\mathrm{v} 4, \mathrm{v} 3, \mathrm{v} 6]$ |

Half-Edge: [v1, v2] or [v2, v1] ?
Should be consistent: e.g. CCW in our configuration

Note*: the container could be array, list, binary search tree...
(it depends, our sample codes used list)


## Half-Edge Data Structure (example)

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| The Edge Container | $[\mathrm{v} 1, \mathrm{v} 3],[\mathrm{v} 1, \mathrm{v} 2],[\mathrm{v} 2, \mathrm{v} 3],[\mathrm{v} 1, \mathrm{v} 4],[\mathrm{v} 3, \mathrm{v} 4], \ldots$ |
| The Face Container | $\mathrm{f} 1[\mathrm{v} 1, \mathrm{v} 2, \mathrm{v} 3] \ldots \mathrm{f} 5 \mathrm{v} 4, \mathrm{v} 3, \mathrm{v} 6]$ |

2). Relationship between primitives:


## Using Half-Edge Data Structure

1. How to check whether a vertex/edge/face is on the boundary?
2. How to track the boundary?
3. How to find your one-ring neighbor?
4. How to do subdivision/simplification...?


## Half-Edge Data Structure (cont.)

## Warm-up Assignment:

Compile and run the "meshlib" codes; use it to load a mesh

1) Go through "iterators.h" and "mesh.h", to see how you can traverse global/local elements.
2) Go through "read()" method, to see how this structure is built up.


# Half-Edge Data Structure (cont.) 

Questions about Half-Edge Data Structure, or the assignment?


## Some 3D Models in Polygonal Meshes

Before we can design a fully robust/powerful GUI and visualization system (which you may keep doing through the semester), here are 3D shapes for you to play a little bit with:
$\square$ Some mesh data (.m format) can be downloaded at: http://www.ece.Isu.edu/xinli/teaching/meshdata1.zip

- A small viewer "G3dOGL.exe" (for .m format mesh) can be downloaded at: http://www.ece.Isu.edu/xinli/Tools/G3dOGL.exe (you can drag your downloaded ".m" file into it directly)
- Many 3D shapes/data online (but in various formats):
- Stanford 3D Scanning Repository: http://graphics.stanford.edu/data/3Dscanrep/
- Aim@Shape Repository: http://shapes.aim-at-shape.net/index.php

Google 3D warehouse

