



Problem 1: Consider the following logic function in canonical form:

$$\sum_{a,b,c,d} m(0, 2, 5, 8, 10, 12, 13).$$

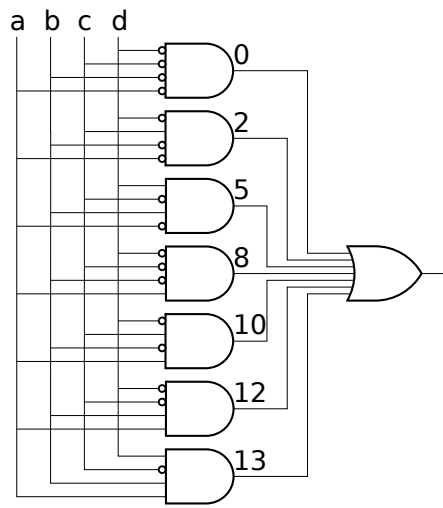
(a) Draw a truth table for this logic function.

SOLUTION:

row	a	b	c	d	!	f
0	0	0	0	0	!	1
1	0	0	0	1	!	0
2	0	0	1	0	!	1
3	0	0	1	1	!	0
4	0	1	0	0	!	0
5	0	1	0	1	!	1
6	0	1	1	0	!	0
7	0	1	1	1	!	0
8	1	0	0	0	!	1
9	1	0	0	1	!	0
10	1	0	1	0	!	1
11	1	0	1	1	!	0
12	1	1	0	0	!	1
13	1	1	0	1	!	1
14	1	1	1	0	!	0
15	1	1	1	1	!	0

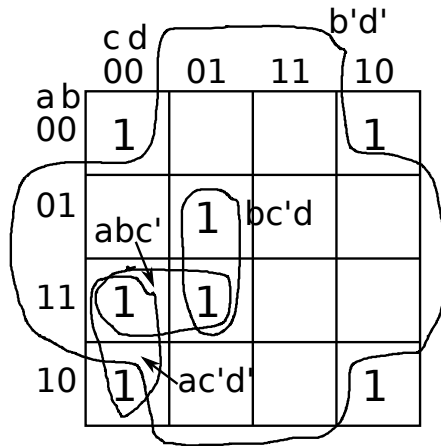
(b) Draw a logic circuit for this function. Do not simplify.

Diagram appears below. The and gates are labeled with minterm numbers.



(c) Draw a Karnaugh map for the logic function.

Solution appears below, the prime implicants are circled.



(d) List the prime implicants.

As shown in the Karnaugh map above, the prime implicants are: $b'd'$, abc' , $bc'd$, and $ac'd'$.

(e) List the essential prime implicants.

The essential prime implicants are $b'd'$ (because of minterms 0, 2, and 10), and $bc'd$ (because of minterm 5).

(f) List all of the minimum cost sum-of-product expressions.

There are only two,

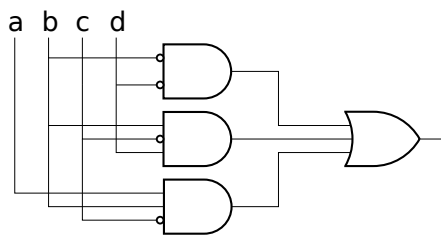
$$b'd' + bc'd + ac'd'$$

and

$$b'd' + bc'd + abc'$$

(g) Draw a logic diagram for your favorite one.

Diagram appears below.

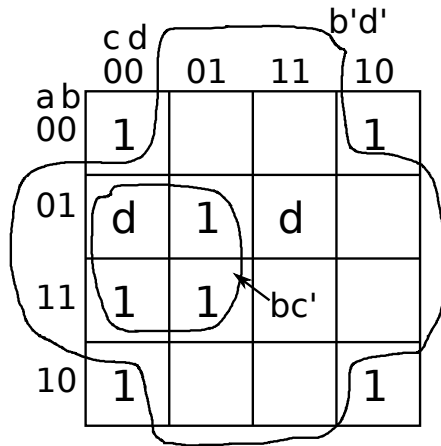


Problem 2: Consider again the logic function from the previous problem,

$\sum_{a,b,c,d} m(0, 2, 5, 8, 10, 12, 13)$. This time however suppose the outputs are *don't care* for two sets of inputs, $a = 0, b = 1, c = 0, d = 0$ (corresponding to row (minterm) 4) and $a = 0, b = 1, c = 1, d = 1$ (corresponding to row (minterm) 7).

(a) Draw a Karnaugh map, include the don't cares.

Solution appears below. The d for minterm 4 (the one on the left) is set to 1, the other is set to zero. This enables the minterms in three "real" non-corner minterms to be covered by just one prime implicant, bc' , which happens to be an essential prime implicant.

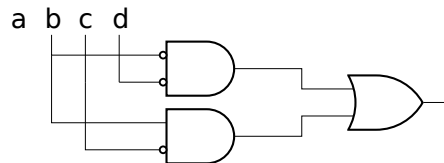


(b) Find a minimum-cost sum-of-products expression making the best use of the don't cares.

Based on the diagram above the two essential prime implicants cover the whole expression. So the minimum cost expression is $b'd' + bc'$.

(c) Draw a logic diagram corresponding to the minimum-cost expression.

Solution appears below.



Problem 3: The *population* of an n -bit quantity is the number of bits with value 1. For example, the population of 4-bit quantity 0101 is 2, the population of 1101011 is 5.

(a) Show a truth table for a Boolean function with an output that's logic 1 if the population of 2-bit input a_1a_0 is the same as the population of 2-bit input b_1b_0 . (The function has four inputs, a_1 , a_0 , b_1 , and b_0 .)

SOLUTION:

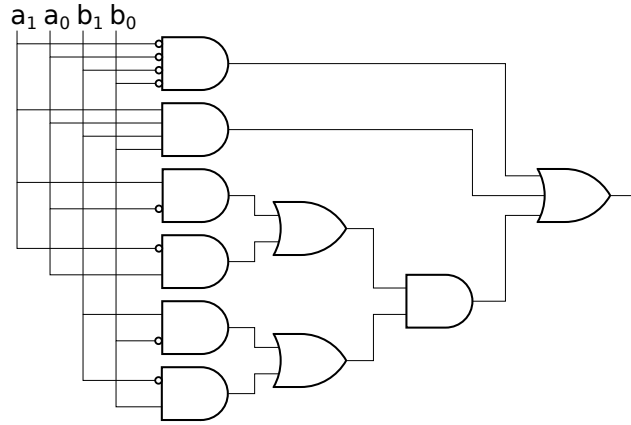
row	a1	a0	b1	b0	!	
0	0	0	0	0	!	1
1	0	0	0	1	!	0
2	0	0	1	0	!	0
3	0	0	1	1	!	0
4	0	1	0	0	!	0
5	0	1	0	1	!	1
6	0	1	1	0	!	1
7	0	1	1	1	!	0
8	1	0	0	0	!	0
9	1	0	0	1	!	1
10	1	0	1	0	!	1
11	1	0	1	1	!	0
12	1	1	0	0	!	0
13	1	1	0	1	!	0
14	1	1	1	0	!	0
15	1	1	1	1	!	1

(b) Derive a Boolean algebraic expression for the same function without using the truth table. Use the following approach: derive an expression that's logic 1 when the population of a_1a_0 is zero. Derive similar expressions for when the population is 1 and when the population is 2. Then pair such expressions for a and b .

The three cases are population 0, 1, and 2. For population 0 both bits of a and b must be 0, that is true when $a'_1a'_0b'_1b'_0$ is true. For population 2 both bits of a and b must be 1, that happens when $a_1a_0b_1b_0$ is true. Population 1 can happen two ways, either $a_1 = 1$ and $a_0 = 0$ or $a_1 = 0$ and $a_0 = 1$. The expression for that for a is $a_1a'_0 + a'_1a_0$. The population for both a and b will be 1 when both expressions are true: $(a_1a'_0 + a'_1a_0)(b_1b'_0 + b'_1b_0)$. Combining all of these conditions yields the solution: $a'_1a'_0b'_1b'_0 + a_1a_0b_1b_0 + (a_1a'_0 + a'_1a_0)(b_1b'_0 + b'_1b_0)$.

(c) Draw a logic diagram for either the hand-derived expression (the previous part) or if you couldn't do the previous part, an expression based on the truth table.

The diagram below is for the hand-derived expression:



(d) Try simplifying the Boolean expressions using the exclusive or (\oplus) operator ($a \oplus b = ab' + a'b$). If successful, draw a logic diagram based on the simplified expressions.

Note that $a_1a'_0 + a'_1a_0 = a_1 \oplus a_0$. Using that we get for the entire expression:

$$a'_1a'_0b'_1b'_0 + a_1a_0b_1b_0 + (a_1 \oplus a_0)(b_1 \oplus b_0).$$

The logic diagram appears below.

