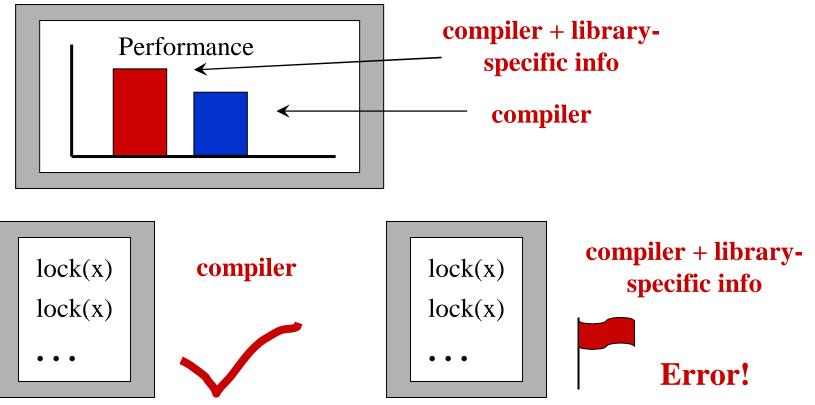
## Compiler Support for Software Libraries

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## Motivation

Numerous libraries exist

• There's a huge benefit to providing compiler support for libraries



# **Outline: Compiler Support for Libraries**

- Requirements
- Our Solution
- Conclusions

#### **Optimization Example**

Consider a dot product routine

```
void dot-product(X,Y,Z){
  for (int i=0,i<n; i++){
    X[i] = Y[i]* Z[i];
  }
}</pre>
```

Consider a common transformation to improve locality

}

2T[i] = t1 \* C[i];

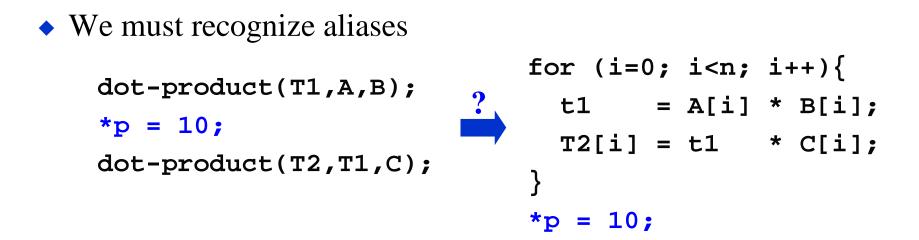
## Syntactic Manipulation is Limited

• We must preserve data dependences

dot-product(T1,A,B);
A[0] = 10;
dot-product(T2,T1,C);

for (i=0; i<n; i++){
 t1 = A[i] \* B[i];
 T2[i] = t1 \* C[i];
}
A[0] = 10;

## Syntactic Manipulation is Limited



• We must correctly handle interactions between the library and the application program

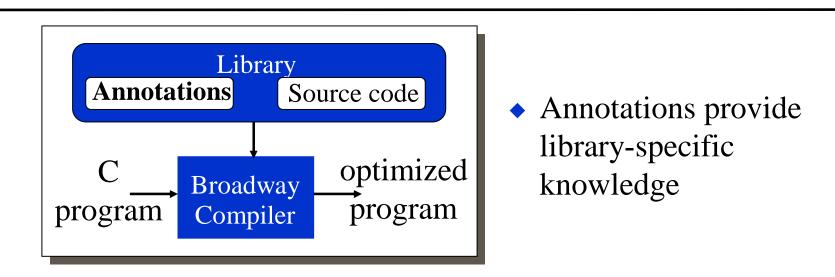
#### What Do We Need?

- Barriers to optimization
  - Data dependences
  - Pointers and aliasing
  - Control flow
  - Complex data structures
- We need the same analyses that traditional compilers use
  - Control flow analysis
  - Data-flow analysis
  - Pointer and dependence analysis

# Are Traditional Compilers Sufficient?

- Libraries are lightweight domain-specific languages
   Compilers need to understand the semantics of these languages
- ◆ Each library has its own semantics
   ⇒ We'd like one compiler for all libraries
- ◆ Each domain specific language is embedded in a base language
   ⇒ We'd like our compiler to understand both languages and the interactions between them

## Our Solution: The Broadway Compiler



- One compiler for all libraries
- Common theme:
  - Expose traditional compiler facilities so that they can be easily configured
  - Integrate the use of these facilities to apply to both libraries and the base language

# **Optimization Opportunities**

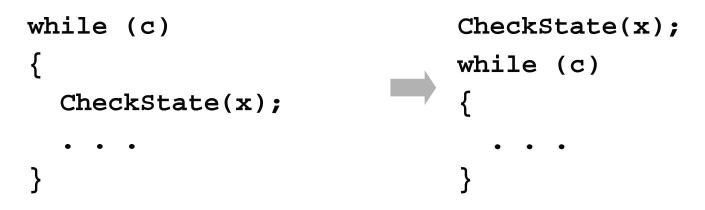
- I. Traditional optimizations on library operators
- **II**. Specializations of library routines
- III. Extensions of traditional optimizations to library operators

requires increasing library-specific information

# I. Traditional Optimizations

#### Trivial example

Loop invariant code motion

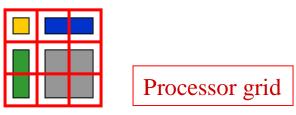


• Requires dependence analysis (or annotations)

# **II. Library Specialization**

Idea

- Analyze dynamic program properties
- Use this information to specialize routines
- Consider a parallel matrix computation
  - Submatrices can have special properties



- Can replace a parallel algorithm with a sequential one
- Requires library-specific data-flow analysis

# **III.** Extensions of Traditional Optimizations

#### Example

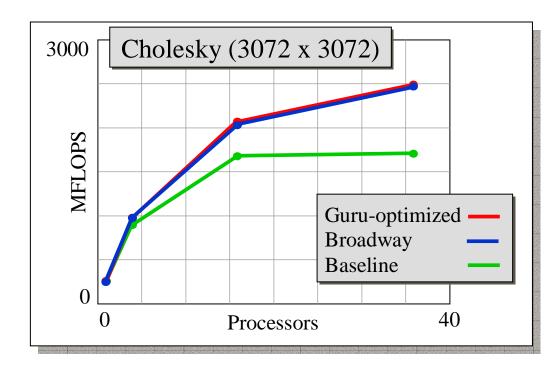
- Constant propagation
- Objects often store state and libraries provide routines to access this state
  - If we can statically determine the state of this object, we can replace function calls with the constant itself
- Requires dependence analysis
- Requires annotations if the state is not stored in an easily accessible form

# The Broadway Compiler

- Two configurable mechanisms
  - Configurable dependence analysis
    - Procedure side effects
    - Pointer relationships
  - Configurable data-flow analysis
- Configurations specified through annotations
- Integrated with built-in mechanisms
  - Aggressive context- and flow-sensitive pointer analysis
  - Various standard optimizations

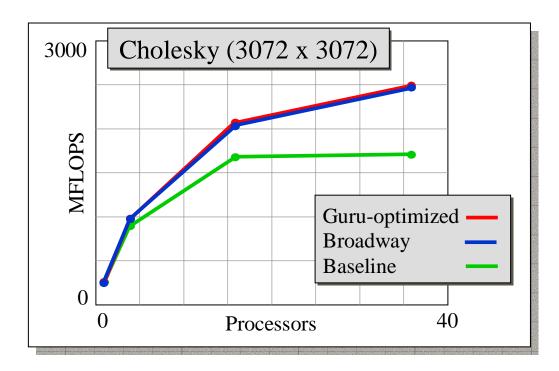
#### Performance Results

- Applied to unmodified PLAPACK parallel dense linear algebra library [van de Geijn 1997]
- Unmodified application and library source code



## **Observations from PLAPACK Results**

- Interactions among multiple optimizations are essential
- Interaction between library and base language are important
- There is benefit to optimizing at multiple levels of abstraction



#### Future Work

- Many other uses of domain-specific compilation
  - Can check for program errors
    - Broadway has been used to identify security holes (Format String Vulnerabilities)
    - More precise than other approaches [Berger, Guyer, Lin 2001]
  - Can remove overhead of language interoperability for PETSc

## Conclusions

- Aggressive program analysis is important
- Significant performance gains possible
- The big picture:

Application level

Library level

Language level

Integrate optimizations across multiple levels of abstraction